

## Standard Notation and Glossary of Terms

**Primitive terms:** players, strategies, payoffs/utilities

### PLAYERS

Individual players	$i = 1, 2, \dots, n$
The set of all players	$I = \{1, 2, \dots, n\}$
Player's $i$ opponents = all the other players except player $i$	$-i$

### STRATEGIES

Specific strategy of player $i$	$s_i$
If we need to enumerate all $k$ strategies of $i$ we write	$s_i^1, s_i^2, \dots, s_i^k$
The set of all strategies (strategy set) of player $i$	$S_i$
Strategy profile is a vector	$s = (s_1, s_2, \dots, s_n)$
Strategy profile of $i$ 's opponents	$s_{-i} = (s_1, s_2, \dots, s_{i-1}, s_{i+1}, \dots, s_n)$
The set of all strategy profiles	$S = S_1 \times S_2 \times \dots \times S_n$

### PAYOFFS

Payoff function of player $i$ is a function (Re denotes the set of real numbers)	$u_i: S \Rightarrow \text{Re}$
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Game,  $G$ , is defined as a triple consisting of the set of players, the set of all strategy profiles and the set of players' utility functions:  $G = \langle I, S, (u_1, u_2, \dots, u_n) \rangle$

## Games in normal form

$p \xrightarrow{\hspace{1cm}} q$  means  $p > q$  . Think about the arrow as: you would go from  $p$  to  $q$ .  
 $p \xleftrightarrow{\hspace{1cm}} q$  means  $p = q$

### The dominance relations, an example

		X		Y		Z
A	1	3		2		2
	↓					
B	3	2		3		0
	↓					
C	2	1		0		1
	↓					
	↓	2		1		1

Strategy B of the Row player **strictly dominates** his strategy A: for each strategy of the Column player, strategy B gives Row a higher payoff.

		X		Y		Z
A	1	3		2		2
	↓					
B	3	2		3		0
	↓					
C	2	1		0		1
	↓					
	↓	3		2		1

Strategy B of the Row player **weakly dominates** his strategy C: for each strategy of the Column player, strategy B gives Row a payoff which is at least as large as the payoff to C and in some cases the payoff is strictly larger.

**Using the general notation to define the two concepts of dominance.**

	$s_2^1$		$s_2^p$		$s_2^m$
$s_1^1$	$u_2(s_2^1, s_1^1)$	.....	$u_2(s_2^p, s_1^1)$	.....	$u_2(s_2^m, s_1^1)$
	$u_1(s_1^1, s_2^1)$	.....	$u_1(s_1^1, s_2^p)$	.....	$u_1(s_1^1, s_2^m)$
.....	.....	.....	.....	.....	.....
$s_1^t$	$u_2(s_2^1, s_1^t)$	.....	$u_2(s_2^p, s_1^t)$	.....	$u_2(s_2^m, s_1^t)$
	$u_1(s_1^t, s_2^1)$	.....	$u_1(s_1^t, s_2^p)$	.....	$u_1(s_1^t, s_2^m)$
.....	.....	.....	.....	.....	.....
$s_1^r$	$u_2(s_2^1, s_1^r)$	.....	$u_2(s_2^p, s_1^r)$	.....	$u_2(s_2^m, s_1^r)$
	$u_1(s_1^r, s_2^1)$	.....	$u_1(s_1^r, s_2^p)$	.....	$u_1(s_1^r, s_2^m)$
.....	.....	.....	.....	.....	.....
$s_1^n$	$u_2(s_2^1, s_1^n)$	.....	$u_2(s_2^p, s_1^n)$	.....	$u_2(s_2^m, s_1^n)$
	$u_1(s_1^n, s_2^1)$	.....	$u_1(s_1^n, s_2^p)$	.....	$u_1(s_1^n, s_2^m)$

In the game above:

$s_1^r$  strictly dominates  $s_1^t$  if for all  $p=1,2,\dots,n$ ,  $u_1(s_1^t, s_2^p) < u_1(s_1^r, s_2^p)$ .

$s_1^r$  weakly dominates  $s_1^t$  if for all  $p=1,2,\dots,n$ ,  $u_1(s_1^t, s_2^p) \leq u_1(s_1^r, s_2^p)$  and for at least one  $p$ ,  $u_1(s_1^t, s_2^p) < u_1(s_1^r, s_2^p)$ .

**THE (GENERAL) DEFINITIONS:**

In any game  $G = \langle I, S, (u_1, u_2, \dots, u_n) \rangle$ , strategy  $s_i$  strictly dominates strategy  $s_i^*$  if for all strategies  $s_{-i}$  of i's opponents  $u_i(s_i^*, s_{-i}) < u_i(s_i, s_{-i})$ ; strategy  $s_i$  weakly dominates strategy  $s_i^*$  if for all strategies  $s_{-i}$  of i's opponents  $u_i(s_i^*, s_{-i}) \leq u_i(s_i, s_{-i})$  and for at least one  $s_{-i}$ ,  $u_i(s_i^*, s_{-i}) < u_i(s_i, s_{-i})$ .

We say that strategy  $s_i$  dominates strategy  $s_i^*$  if  $s_i$  dominates  $s_i^*$  either strictly or weakly; such  $s_i^*$  will be called dominated.

We say that strategy  $s_i$  is dominant if it dominates all the other strategies of i.

Strategy  $s_i$  is the best response to  $s_{-i}$  if  $u_i(s_i, s_{-i}) \geq u_i(s_i^*, s_{-i})$  for all strategies  $s_i^*$  of player I. Denote:  $s_i = BR_i(s_{-i})$

A profile of strategies  $(s_1, s_2, \dots, s_n)$  is in Nash equilibrium if for every  $i=1,\dots,n$ ,  $s_i = BR_i(s_{-i})$ .

